Construction of linear involutory transformations over finite fields through the multiplication of polynomials modulus a polynomial

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Introduction

Let $\mathcal{P} = \mathbb{F}_q$ be a finite field of $q = p^t$ elements being p a prime number and $t \in \mathbb{N}$. We have that $\mathcal{P}[x]$ is the ring of polynomials having coefficients in \mathcal{P} and using the polynomial $F(x) \in \mathcal{P}[x]$ let us build the ring $\mathcal{R} = \mathcal{P}[x]/_{F(x)}$. Let $g(x), \mu(x)$ be elements of \mathcal{R} we define the transformation $\Psi : \mathcal{R} \to \mathcal{R}$ as

$$\Psi(g(x)) = g(x)\mu(x) \mod F(x).$$

Our Interest

- \bullet Conditions Conditions for the Ψ transformation to be involutory?
- \bullet Construction MDS involutory matrices from Ψ

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Relation between nilpotent and involutory elements

Definition

• The element $r \in \mathcal{R}$, $r \neq 0$ is called a zero divisor if exists $r' \in \mathcal{R}$, $r' \neq 0$ such that

$$r * r' = 0.$$

2 A zero divisor $r \in \mathcal{R}$ is called nilpotent if exists $n \in \mathbb{N}$ such that

$$r^n = 0.$$

The lowest n = n(r) which satisfies the previous property is called nilpotency index.

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- 1. If \mathcal{R} contains at least one involutory element $r \in \mathcal{R}$ then \mathcal{R} contains at least one nilpotent element with nilpotency index 2.
- 2. Let r' be one involutory element of \mathcal{R} then the remaining involutory elements r are uniquely defined by

$$r = r' + w$$

where w is a nilpotent element with nilpotency index 2, i.e., any two involutory elements are related by a nilpotent element.

Let denote D(2) the union of all the ideal $\mathcal{R}a$ where a is a nilpotent element with nilpotency index 2 and let $b \in D(2)$ be a fixed element, then for any $r \in \mathcal{R}$ such that $rb \neq 0$ the element

$$a = e + rb \tag{2}$$

is an involutory element. Hence, any transformation $\Psi:\mathcal{R}\to\mathcal{R}$ constructed as

$$\Psi(x) = x(e+rb) \tag{3}$$

is a linear involutory transformation.

Let be \mathcal{P} finite field and $\mathcal{R} = P[x]/_{F(x)}$ where $F(x) \in P[x]$, the transformation $\Psi : \mathcal{R} \to \mathcal{R}$ is defined as

$$\Psi(g(x)) = g(x)\mu(x) \mod F(x).$$

considering that F(x) has canonical factorization

$$F(x) = f_1(x)^{k_1} \cdots f_s(x)^{k_s}$$

with $\deg(F(x)) = m$ and $\deg(f_i(x)) = m_i$, such that $m = k_1 m_1 + \dots + k_s m_s$.

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The amount of polynomials $h(x) \in \mathcal{P}[x]$ such that $\deg(h(x)) < \deg(F(x))$ and $\gcd(h(x), F(x)) = 1$ is

$$q^m \left(1 - \frac{1}{q^{m_1}}\right) \cdots \left(1 - \frac{1}{q^{m_t}}\right) \tag{4}$$

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Theorem

The ring \mathcal{R} contains involutory elements iff one of the following conditions holds

- 1. p is an odd value
- 2. If p = 2 in the canonical factorization of F(x) exists $i \in \{1, \ldots, t\}$ such that $k_i \ge 2$.

Corollary 1

For any $n \in \mathbb{N}$, n > 1 it is possible to build a ring \mathcal{R} of p^{tn} elements which contain involutory elements.

Corollary 2

For any $n \in \mathbb{N}$ such that n > 1 exist one linear involutory transformation

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$$\Psi: \mathcal{R} \to \mathcal{R}$$

build by the rule in (1).

If the order of F(x), denoted by e, is even, then the application $\Psi : \mathcal{R} \to \mathcal{R}$ of the form

$$\Psi(g(x)) = g(x)x^{\frac{e}{2}} \mod F(x)$$
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is involutory.

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Branch Number

Let be function $\tau : \mathcal{R} \to \mathbb{N}_0$ defined as:

$$\tau(g(x)) = \sum_{i=0}^{m-1} \delta_{0,f_i}$$

where

$$\delta_{0,f_i} = \begin{cases} 1, \text{ if } f_i \neq 0\\ 0, \text{ if } f_i = 0, \end{cases}$$

Definition

The branch number ρ of a linear transformation $\Psi : \mathcal{R} \to \mathcal{R}$ is defined as

$$\rho(\Psi) = \min_{g(x)\neq 0} \{\tau(g(x)) + \tau(\Psi(g(x)))\}.$$

The linear transformation Ψ is called MDS if $\rho(\Psi) = m + 1$.

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Let $(\alpha_0, \ldots, \alpha_s)$ be elements of the set \mathcal{P}^* , such that $\forall i, j \in \{0, \ldots, s\} \; \alpha_i \neq \alpha_j$ and let be the polynomials such that

$$f(x) = (x + \alpha_0)^{n-1} (x + \alpha_1) \cdots (x + \alpha_s)$$
 and $F(x) = (x + \alpha_0) f(x)$.

For all element $h(x) \in \mathcal{R}$ the transformation Ψ defined by the rule (1) is a linear involutory transformation whose branch number satisfies that

$$\rho(\Psi) \leqslant 4.$$

If the matrix associated to Ψ does not contain any zero coefficient then

$$\rho(\Psi) = 4.$$

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Let a be element of field \mathcal{P} such that $a \neq 1$, and let be the polynomials such that

$$f(x) = (x^n + a)$$
 and $F(x) = f(x)^2$.

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Example MDS involutory transformation

Take $\mathcal{P} = \mathbb{F}_2[x]/_{x^8+x^4+x^3+x^2+1}$ having β a primitive element on \mathcal{P} and selecting the following parameters

n	f(x)	F(x)	h(x)
4	$(x+\beta)(x+\beta^2)$	$f(x)^2$	x
6	$(x+\beta)(x+\beta^2)(x+\beta^3)$	$f(x)^2$	$x^2 + x$

the transformation

$$\Psi(g(x)) = g(x)(1 + f(x)h(x)) \mod F(x)$$

is an involutory MDS transformation.

For the class of application Ψ such that

$$\Psi(g(x)) = g(x)(1 + (x^n + a)h(x)) \mod x^{2n} + a^2,$$

 $\forall a \in \mathcal{P}^* \ y \ \forall \ h(x) \in \mathcal{R}$ does not exist MDS matrices.

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Generation of 4×4 involutory MDS

Let $\mathcal{P} = \mathbb{F}_{2^t}$, given $a \in \mathcal{P}^*$, $f_1(x) = x^n + a$, $F_1(x) = x^{2n} + a^2$, $f_2(x) = x^n + 1$, $F_2(x) = x^{2n} + 1$, we construct the transformations Ψ_1 and Ψ_2 from \mathcal{P}_i to \mathcal{P}_i where $\mathcal{P}_i = P[x]/_{F_i(x)}$, i = 1, 2

$$\Psi_1(g(x)) = g(x)(1 + f_1(x)h(x)) \mod F_1(x)$$

and

$$\Psi_2(g(x)) = g(x)(1 + f_2(x)x) \mod F_2(x).$$

Let us denote by " \circ " the composition of two applications and by $\mathcal{R}(2n)$ the set of all polynomials with coefficients in \mathcal{P} whose degree is lower than 2n. We define the transformation

$$\Psi: \mathcal{R}(2n) \to \mathcal{R}(2n)$$

of the following form

 $\forall g(x) \in \mathcal{R}(2n), \ \Psi(g(x)) = \Psi_1 \circ \Psi_2 \circ \Psi_1(g(x)). \tag{6}$ (CTCRYPT'2023) Institute of Cryptography, Cuba June 6, 2023 15 / 20

Let be n = 4. For a given polynomial $h(x) = h_3 x^3 + \cdots + h_0$ such that $h_1 \neq ah_3$ and for any value of h_2 there exist a value of h_0 such that the matrix associated to transformation Ψ is MDS. Particularly one can take $h_1 = ah_3 + 1$.

Algorithm

INPUT: $a \in \mathcal{P} \setminus \{0\}$ and $h_2, h_3 \in \mathcal{P}$

STEP 1: Make $h_1 = ah_3 + 1$, i = 1.

STEP 2: Make the i-th row of the matrix \mathcal{A}_{h_0} , $\delta_i(x) = \Psi(x^{i-1})$.

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Algorithm

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STEP 3.1: Discard all roots of the polynomials $\delta_{i,1}(h_0), \ldots, \delta_{i,4}(h_0)$. STEP 3.2: Make i = i + 1

STEP 3.3: If i = 5 go to STEP 4, otherwise go to STEP 2.

- STEP 4: Discard from \mathcal{A}_{h_0} all the roots of the polynomials result from calculating its order 2 minors.
- **RETURN:** The set of values for h_0 such that the polynomials $h(x) = h_3 x^3 + h_2 x^2 + h_1 x + h_0$ generate an involutory MDS matrix.

Make $\mathcal{P} = \mathbb{F}_2[x]/_{x^8+x^4+x^3+x^2+1}$ having β a primitive element on \mathcal{P} and selecting the following parameters:

- $f(x) = x^2 + \beta$
- $F(x) = f(x)^2$
- $h_3 = 1$
- $h_2 = 0$

Applying the above procedure we obtain 230 distinct values of h_0 such that the polynomial $h(x) = h_3 x^3 + h_2 x^2 + h_1 x + h_0$ generate an involutory MDS matrix.

The method presented in this section to generate involutory MDS matrices can be also used to generate key-dependent MDS matrices for dynamic encryption algorithms, for example in following work authors using this point of view

Pablo Freyre, Oristela Cuellar, Nelson Dáz, Adrián Alfonso. From AES to Dynamic AES. Journal of Science and Technology on Information Security, 2020.

Pablo Freyre, Oristela Cuellar, Nelson Díaz, Ramses Rodríguez, Adrián Alfonso. Dynamic Cryptographic Algorithms Kuznyechik and Magma. Journal of Science and Technology on Information Security, 2020.

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Conclusion

We board the construction of involutory elements using the nilpotent elements of a ring with nilpotency index 2. From such construction we study the involutory linear applications generated by multiplication of polynomials with coefficients in a finite field, characterizing the branch number of such applications. Finally, we introduce a new proposal which allow to generate MDS matrices using the results of the paper towards involutory elements and alternating between two different rings. However, how to reduce from quadratic to linear equations, which allow to increase the dimension of the matrix, remains as open question and future line of work in our research.