



Security Code

On Efficiency of Block Encryption by Improved Key Schedule

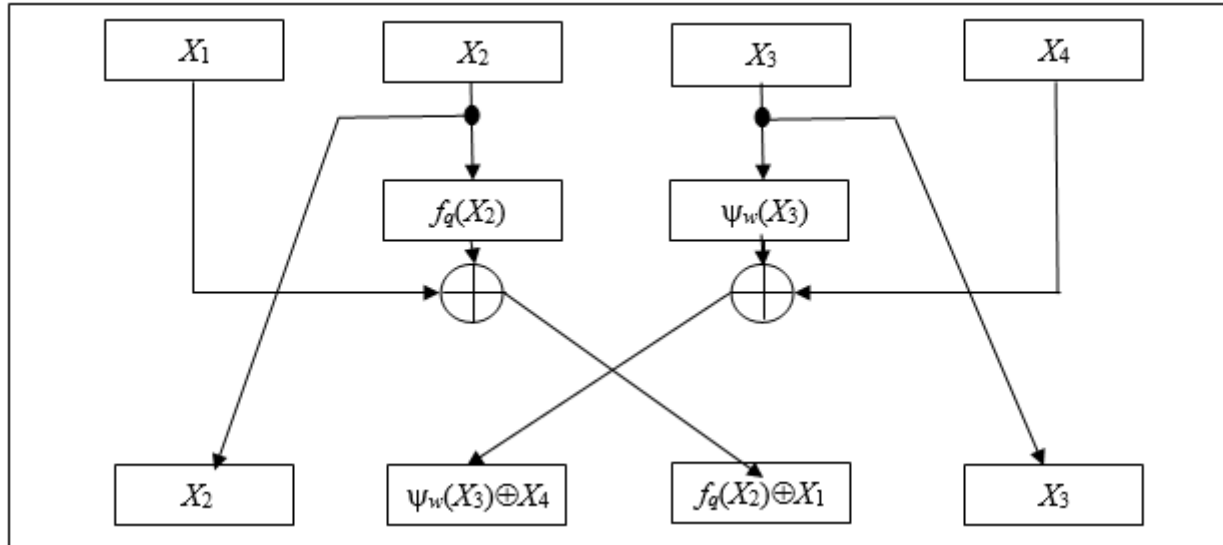
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Introduction

- Approach to construction of block encryption algorithm with higher efficiency compared with GOST 28147-89;
- Increasing of efficiency by 2 – 2,6 times by reducing number of encryption rounds to 12 – 16;
- Improved key schedule for algorithm based on generalized Feistel network with using GOST 28147-89 combining function.
- Estimation of encryption speed must be made more precise by modern software tools.

Block encryption network based on shift register of length 4



Encryption: $T g_{q_r, w_r} \cdots g_{q_2, w_2} g_{q_1, w_1} (X_1, X_2, X_3, X_4) = (Y_1, Y_2, Y_3, Y_4)$, q_i, w_i – keys for round $i=1, \dots, r$

Decryption: $T g_{q_1, w_1} T^2 \cdots T^2 g_{q_{r-1}, w_{r-1}} T^2 g_{q_r, w_r} (Y_1, Y_2, Y_3, Y_4) = (X_1, X_2, X_3, X_4)$.

Permutation T : $T(Z_1, Z_2, Z_3, Z_4) = (Z_3, Z_1, Z_4, Z_2)$.

Round permutation $g_{q,w}$: $g_{q,w}(X_1, X_2, X_3, X_4) = (X_2, \psi_w(X_3) \oplus X_4, f_q(X_2) \oplus X_1, X_3)$.

Combine functions f_q and ψ_w are defined on the next slide.

Mixing properties of round permutation $g_{q,w}$

Combine functions: $f_q = T_z(S_{8,4}(X_2 \oplus q))$; $\psi_w = T_z(S_{8,4}(X_3 \oplus w))$,

\oplus – addition modulo 2^{32} ;

T_z – left 11 bits cyclic shift for $z=left$;

T_z – right 11 bits cyclic shift for $z=right$;

$S_{8,4} = (S_1, \dots, S_8)$ – s -boxes of GOST 28147-89;

M – mixing matrix of round permutation $g_{q,w}$;

M' – mixing matrix of permutation $(g_{q,w})^{-1}$;

$\exp M$ ($\exp M'$) – exponents of M (of M');

r – recommended lower bound for number of rounds.

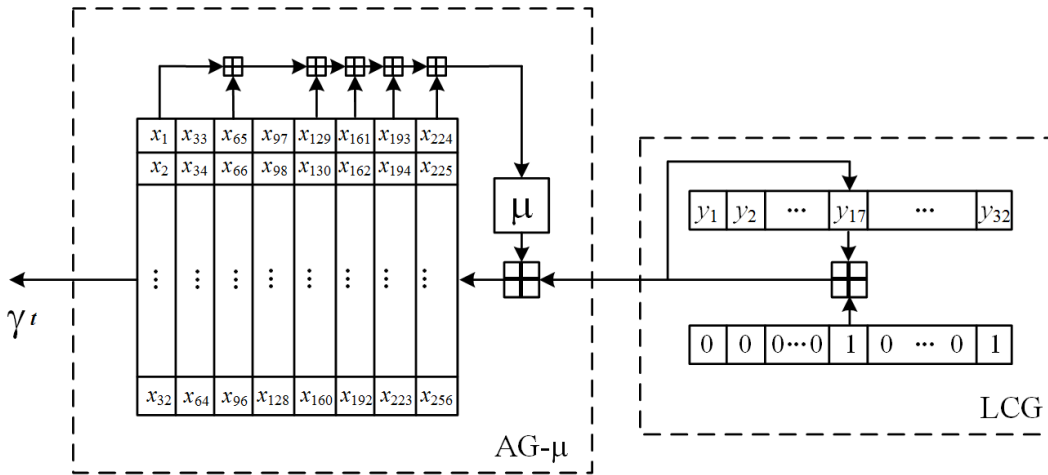
Table 1 – Values of $\exp M$, $\exp M'$, r

	f_q for $z=left$	f_q for $z=right$
ψ_w for $z=left$	8, 8, 15	7, 7, 13
ψ_w for $z=right$	7, 7, 13	7, 6, 12

Note: To counter meet-in-the-middle attack it is enough that for any $l \in \{1, \dots, r-1\}$ full mixing will be achieved after l encryption rounds or after $r-l$ decryption rounds. Hence, number of encryption rounds must be at least 12.

Generator G of round keys

Generator G :



Transformation of G states:

$$(X_{0,t+1}, X_{1,t+1}, \dots, X_{7,t+1}) = (X_{1,t}, \dots, X_{7,t}, f^{\mu}(X_{0,t}, X_{1,t}, \dots, X_{7,t}) \oplus Y_{t+1}),$$

$$f^{\mu}(X_{0,t}, X_{1,t}, \dots, X_{7,t}) = \mu\left(\left(\sum_{j=0}^7 a_j X_j\right) \bmod 2^{32}\right),$$

$$a_1, \dots, a_7 \in \{0, 1\}, a_1 + \dots + a_7 > 0, a_0 = 1.$$

Properties of sequence $\{\gamma_t\} = b_{32}(X_{0,t})$, $t=0, 1, \dots, b_{32}(X) - 1$ – binary 32-bit notation for $X \in \mathbb{Z}_m$, $m=2^{32}$:

- Length of the sequence $\{\gamma_t\}$ period is a multiple of 2^{32} .
- For γ_t all 256 bits of initial state are sufficient for choose a_0, \dots, a_7 and certain transformation μ , $t \geq 16$.

Encryption key: initial state of G ;

Encryption key length: $128 + 32k$ bits, $k \in \{0, 1, 2, 3\}$, integers $X_{0,0}, \dots, X_{3-k,0}$ equals 0.

Set of round keys: determinate non-regular selection from $\{\gamma_t\}$, $t < 100$.

Recommended parameters and Advantages

Recommended parameters

- Size of input block: 128 bits
- Combining functions:
 $f_q(X_2) = T_{\text{right}}(S_{8,4}(X_2 \oplus q))$ and
 $\psi_w = T_{\text{right}}(S_{8,4}(X_3 \oplus w))$,
 S_1, \dots, S_8 – s-boxes of GOST 28147-89.
- Order of key set: $2^{128+32k}$, $k \in \{0, 1, 2, 3, 4\}$.
- Number of encryption rounds: from 12 to 16 with corresponding key schedule.

Advantages

- 2-2,6 times reduction of encryption rounds number (compared with GOST 28147-89) by improved key schedule.
- Ability to set encryption key length is equal to $128+32k$ bits, $0 \leq k \leq 4$.
- Increased block size (compared with GOST 28147-89) with ability to use combine functions of GOST 28147-89.
- A slight difference between encryption and decryption algorithms.

Thank you!