



# Main directions of research and pedagogical activities of M.M. Glukhov's scientific school

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- ▶ Mikhail Mikhailovich Glukhov (1930 - 2018) - Doctor of Physical and Mathematical Sciences, Professor, full member of the Russian Academy of Cryptography, Honored Scientist of the Russian Federation.

# A brief biography

- ▶ Mikhail Glukhov was born on November 20 in 1930 in the Tatar Autonomous Soviet Socialist Republic, in a working-class family.
- ▶ In 1949 he graduated from school and enrolled in the Melekes Teachers' Institute and in 1951 he graduated with honors.
- ▶ In 1951 - 1954 he worked as a mathematics and physics teacher at school and studied at the Ulyanovsk Pedagogical Institute.
- ▶ In 1954 he graduated with honors and started teaching in the Ulyanovsk Pedagogical Institute.
- ▶ In 1958 Mikhail Glukhov as a full-time postgraduate joined Lenin MSPU and continued teaching at the Moscow Correspondence Pedagogical Institute.

# A brief biography

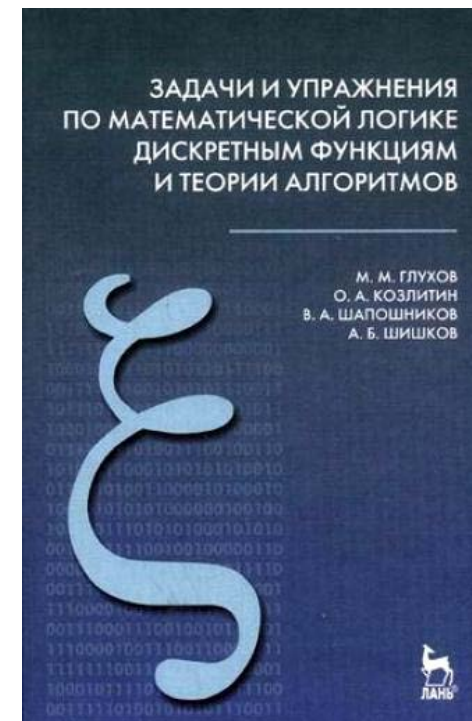
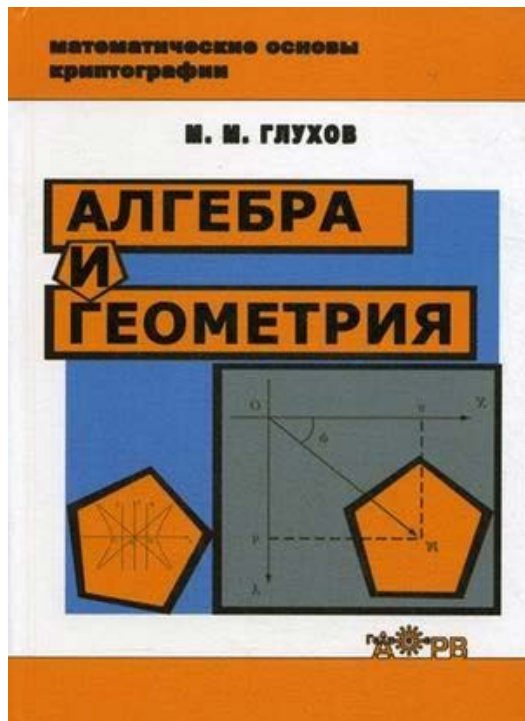
- ▶ At the end of graduate school he became the dean of the Faculty of Physics and Mathematics at Melekes Pedagogical Institute.
- ▶ In 1962 he defended his PhD thesis and then his doctoral dissertation in 1974.
- ▶ In 1973 - 1993 - Director of the Department of Mathematics.
- ▶ In 1993 - 2012 - Professor of the Department of Mathematics.
- ▶ Since 1992 Mikhail Glukhov was a full member of the Academy of Cryptography of the Russian Federation.
- ▶ Since 2005 he was Academician-Secretary of the Department of Mathematical Methods of Cryptography of Academy of Cryptography of the Russian Federation.

# Pedagogical activity

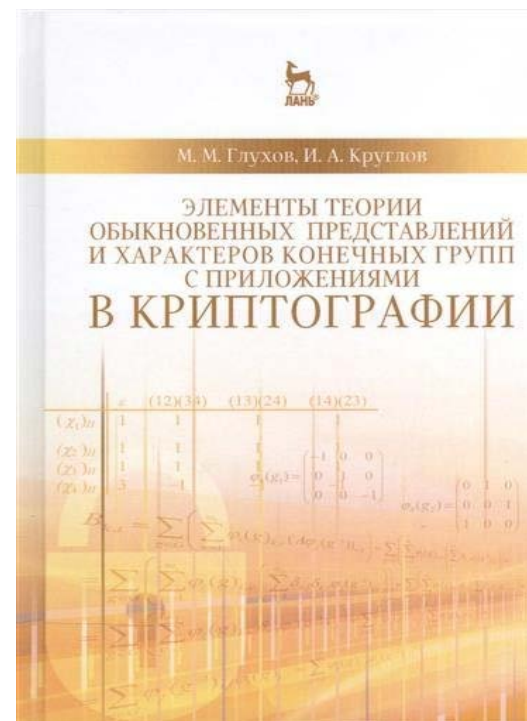
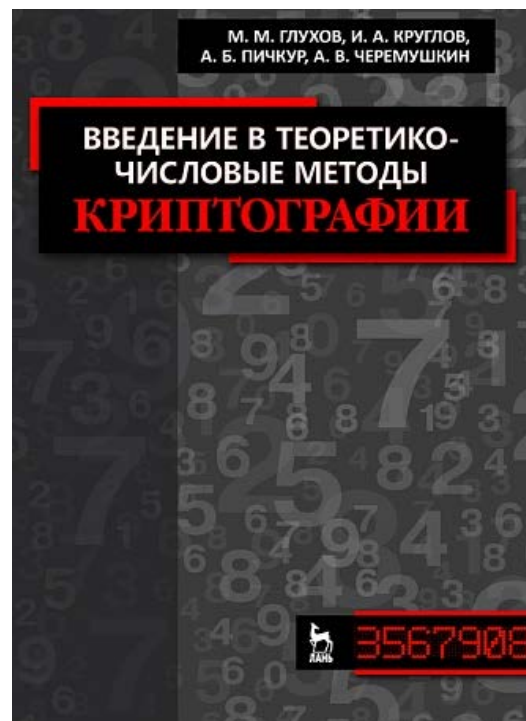
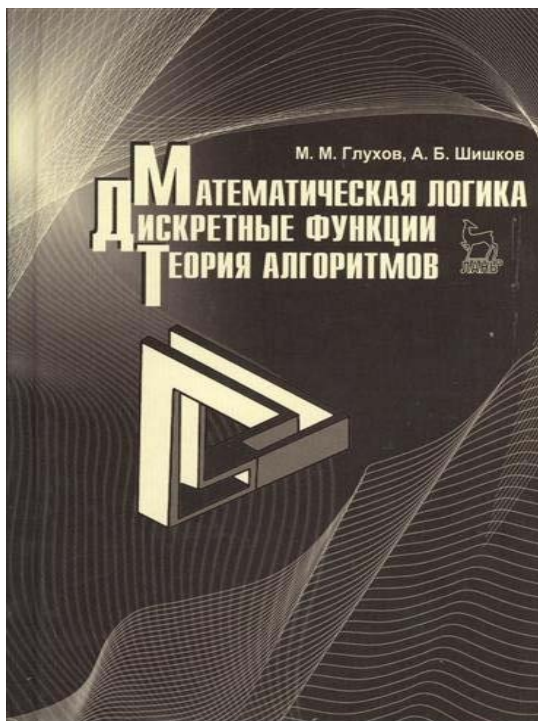
## Textbooks and tutorials:

- ▶ the finite automata theory
- ▶ mathematical logic and theory of algorithms
- ▶ algebra and analytical geometry
- ▶ number-theoretic methods of cryptography
- ▶ group representation theory

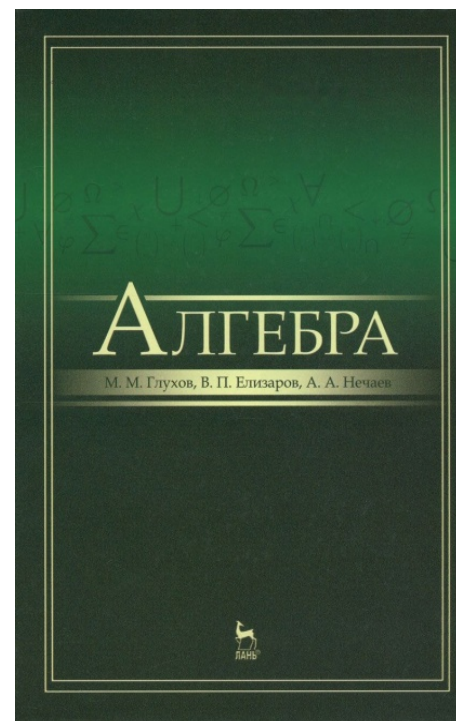
# Pedagogical activity



# Pedagogical activity



# Pedagogical activity





# Algebraic aspects of cryptography

- ▶ At the turn of 1990s – 2000s there was a rapid development of the quantum computation theory
- ▶ In 1997 polynomiality of the discrete logarithm problem was proved in quantum setting
- ▶ This result initiated research in finding other difficult mathematical problems which might be used to design asymmetric cryptosystems

# Algebraic aspects of cryptography

- ▶ 2007 - **Scherbakov V.A., Kostina A.A., Moldovyan P.A., Moldovyan N.A.** Proposed a public-key scheme based on the composition of conjugacy and discrete logarithm problems as a basis for national standard
- ▶ The authors' assumption was that the problem of key recovery in the scheme was not reduced to the discrete logarithm problem in a cyclic group of a finite field
- ▶ Mikhail Glukhov found a vulnerability and put the final point in the discussion about standardization of this scheme
- ▶ The Glukhov's report on this topic was the first talk presented at the 1<sup>st</sup> CTCrypt workshop in 2012

# Algebraic aspects of cryptography

Let

$$A(\varepsilon) = \{(a, b, c) : a, b, c \in \mathbb{Z} / p^2\}$$

over  $\mathbb{Z} / p^2$  ( $p$  is large prime) with coordinate-wise addition of vectors and multiplication of vectors

$$(a, b, c)(x, y, z) = (ax + \varepsilon bz + \varepsilon cy, ay + bx + cz, az + \varepsilon by + cx),$$

where parameter  $\varepsilon \in \mathbb{Z}$ ,  $p \mid \varepsilon$ ,  $0 < \varepsilon < p^2$ .

Identity of algebra  $A(\varepsilon)$  is the vector  $(1, 0, 0)$ , element  $\alpha = (a, b, c)$  is invertible if and only if  $(a, p) = 1$ ,  $\Gamma = A(\varepsilon)^*$ , the order of  $\Gamma$  is equal to  $p^5(p-1)$ .

Based on these data, they concluded that in the general case the group  $\Gamma$  is not cyclic and the maximal order of its cyclic subgroups is equal to  $p^2(p-1)$ .

# Algebraic aspects of cryptography

In the proposed signature scheme the secret key is the pair of elements  $X_1, X_2 \in \Gamma$  of order  $\omega \geq p^2$ , and the public key is the pair  $Y_1 = X_1^p, Y_2 = X_2^p$

The security of the scheme is determined by the complexity of finding a  $p$ -th root of an element of  $\Gamma$ .

The authors considered two algorithms for solving this problem, the fastest of them has the complexity  $O(p\sqrt{p})$  operations in  $\Gamma$ .

As a result they claimed that the group  $\Gamma$  is promising for designing signature schemes.

# Algebraic aspects of cryptography

## Structure of the group $\Gamma$

**Lemma 1.** The set  $H$  of elements  $\alpha = (a, b, c) \in \Gamma$ , where  $p \mid c$ , is a subgroup of  $\Gamma$ ,  $|H| = p^4(p-1)$ , and for all  $\alpha = (a, b, c) \in H$  and  $k \in \mathbb{N}, k \geq 2$  the following equality holds:

$$\alpha^k = (a^k, ka^{k-1}b, ka^{k-2}(ac + C_k^2 \varepsilon b^2)).$$

**Lemma 2.** The set  $K$  of elements  $\alpha = (a, b, c) \in \Gamma$ , where  $p \mid b, p \mid c$ , is a subgroup of  $\Gamma$ ,  $|K| = p^3(p-1)$ , and for all  $\alpha = (a, b, c) \in K$  and  $k \in \mathbb{N}$  the following equality holds:

$$\alpha^k = (a^k, ka^{k-1}b, ka^{k-1}c).$$

# Algebraic aspects of cryptography

## Structure of the group $\Gamma$

**Lemma 3.** Let  $L$  be the set of elements  $\alpha = (a, b, c) \in \Gamma$ , where  $p \mid b$ . Then for all  $\alpha = (a, b, c) \in L$  and  $k \in \mathbb{N}, k \geq 5$  the following equality holds:

$$\alpha^k = (a^k + \varepsilon C_k^3 a^{k-3} c^3, ka^{k-1}b + C_k^2 a^{k-2} c^2 + \varepsilon C_k^5 a^{k-5} c^5, ka^{k-1}c + \varepsilon C_k^4 a^{k-4} c^4).$$

**Lemma 4.** Each vector  $\alpha = (a, b, c)$  from the set

$$M = \{(a, b, c) : (b, p) = 1, (c, p) = 1\}$$

can be represented in the form  $\alpha = \alpha_1 \alpha_2$ , where  $\alpha_1 \in H \setminus K, \alpha_2 \in L \setminus K$ .

# Algebraic aspects of cryptography

## Structure of the group $\Gamma$

**Theorem 1.** The following decompositions of the groups  $K$ ,  $H$  and  $\Gamma$  into the direct product of cyclic groups hold:

$$K = G_p^{(1)} \times G_p^{(2)} \times G_p^{(3)} \times H_{p-1},$$

$$H = G_{p^2}^{(1)} \times G_p^{(2)} \times G_p^{(3)} \times H_{p-1},$$

$$\Gamma = G_{p^2}^{(1)} \times G_{p^2}^{(2)} \times G_p^{(3)} \times H_{p-1},$$

where  $G_{p^k}^{(i)}$  is the group of order  $p^k$ ,  $k = 1, 2, i = 1, 2, 3$ , and  $H_{p-1}$  is the group of order  $p - 1$ .

**Theorem 2.** The complexity of finding the  $p$ -th root of an element  $\xi \in \Gamma$ , is  $O(1)$ , if  $\text{Ord } \xi < p^2$ , and  $O(\log^2 p)$  or  $O(\log^3 p)$ , if  $\text{Ord } \xi \geq p^2$ .

**Corollary.** The complexity of finding the secret key in the examined signature scheme is  $O(\log^3 p)$  binary operations.

# Algebra, cryptography and education

- ▶ A conference dedicated to the problems of teaching algebra in high school specialties in the field of information security was held in 2004
- ▶ The discussion was about which areas of algebra should be included in the student's education
- ▶ Mikhail Glukhov explained his viewpoint on the example of AES



# Algebra, cryptography ...

Textblocks, keys and states are interpreted as matrices over finite field

$$GF(2^8) = GF(2)[x]/p(x), \quad p(x) = x^8 + x^4 + x^3 + x + 1.$$

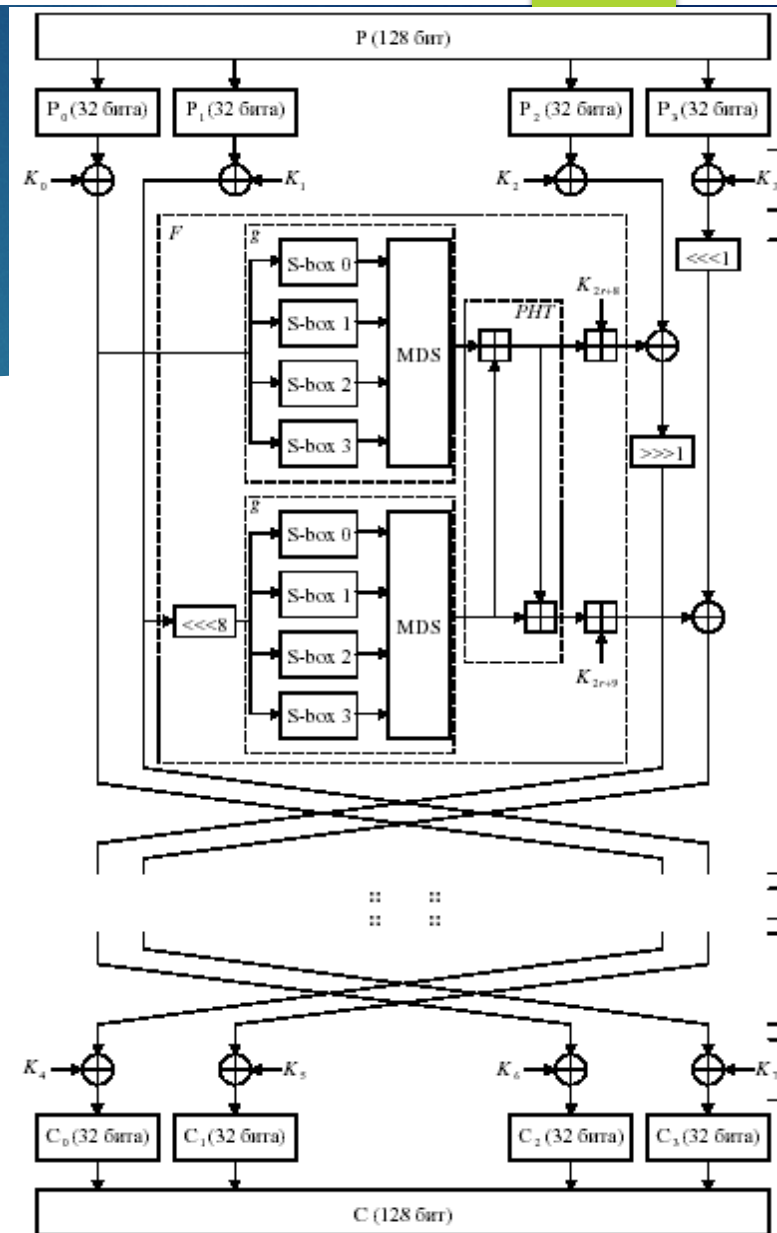
Byte  $b = (b_7, b_6, \dots, b_0)$  is interpreted as the polynomial

$$b(x) = b_7x^7 + b_6x^6 + \dots + b_0 \in GF(2)[x]/p(x).$$

Matrices, which represent textblocks and keys, have 4 rows and 4 columns.

Round transformations have 3 types of slices:

- modulo 2 addition of an input block with an iterative key;
- non-linear transformation,
- linear transformation.



# Algebra, cryptography and education

Non-linear transformation is represented as parallel application of 16 fixed bijective substitutions:  $s_1 = \tau \cdot l \in S(GF(2^8))$ ,  $\tau$  - reverse of non-zero elements of the field  $GF(2^8)$ ,  $l$  - affine transformation of  $GF(2^8)_{GF(2)}$ .

Linear transformation is a composition of maps  $h_1, h_2 \in GL_{16}(2^8)$ :

$h_1$  - byte-wise cyclic shift of the three last rows of a matrix,

$h_2$  - multiplication on circulant matrix, may be represented as a polynomial multiplication modulo polynomial  $x^8 + 1$ .

# Algebra, cryptography and education

Operations in AES:

- linear transformations of the vector space over the field  $GF(2)$ ;
- substitutions;
- finite field  $GF(2^8)$ , as an extension of the field  $GF(2)$  by the root of irreducible polynomial .

# Algebra, cryptography and education

Therefore an information security specialist should know:

- ▶ basics of linear algebra, the theory of vector spaces and linear transformations
- ▶ group theory fundamentals, in particular, the permutation groups theory
- ▶ basics of the theory of finite fields and polynomials over them

Actually, these areas form the basis of the algebra textbook

# Publications and research areas

More than 150 papers in the

- ▶ theory of discrete functions
- ▶ theory of the finite rings
- ▶ theory of quasi-groups and their applications

**The main contribution** – in the theory of quasi-groups, groups and loops (non-associative quasi-groups) and their applications in cryptography

Researches on the coverage length, width and depth of the finite groups:

- ▶ length of coverage and powers of coverage layers for the main systems of generators of the symmetric group
- ▶ the value of the index measure of transitivity of the sets of round transformations of AES

# Publications and research areas

- ▶ Researches in linear recursive sequences over finite rings and their applications
- ▶ Introduced iterative discrete functions and explored their properties
- ▶ Obtained the generalization of A.A. Markov theorem about non-distortion mappings
- ▶ Introduced index of affinity of discrete function, studied its connection with the function spectrum

# Publications and research areas

- ▶ Obtained the group theory criteria of the planarity of the map. The map  $f: GF(2^n) \rightarrow GF(2^n)$  is planar if the map  $f(x+a)+f(x)+ax$  is bijective for all non-zero  $a$ .
- ▶ **Definition.** Let  $G_1, G_2$  - equivalent abelian groups,  $H_1$  and  $H_2$  – their right regular representations. The map  $f: G_1 \rightarrow G_2$  is *planar*, if the map  $f(x+a)-f(x) -$  is bijective for all non-zero  $a$ .
- ▶ **Theorem.** The map  $f$  is *planar* if and only if the set of substitutions  $H_1 f H_2$  is 2-transitive.

# Scientific school

- ▶ Under Mikhail Glukhov guidance 14 PhD theses was defended, three of his students became doctors of science
- ▶ His scientific school was supported by the grant from President of the Russian Federation





# Scientific school

The main directions of research:

- ▶ algebraic, number-theoretic and combinatorial problems of cryptography
- ▶ finite rings and modules theory, the study of identities and corresponding polynomial functions
- ▶ multilinear recurrences and linear codes over finite bimodules
- ▶ perspective methods for constructing pseudo-random sequences with guaranteed cryptographic properties

# Scientific school

The main directions of research:

- ▶ finite groups, permutation groups and corresponding operations with applications to design and analysis of cryptographic mechanisms
- ▶ properties of discrete functions, the construction of functions with specified properties, their classification by cryptographic parameters
- ▶ methods for solving systems of polynomial equations over finite rings using standard bases
- ▶ number-theoretic problems of cryptography

# Scientific school

From 2003 to 2013 - 429 publications, including:

- ▶ 3 monographs
- ▶ 22 textbooks and manuals
- ▶ 195 articles in leading scientific journals
- ▶ 209 talks, of which 35 publications in foreign scientific journals

Defended: 6 doctoral and 15 PhD theses

# Awards

Mikhail Glukhov was awarded:

- ▶ the Order of the Badge of Honor in 1986
- ▶ the Order of Honor in 2010 for success in teaching and research
- ▶ the title of Honored Scientist of the Russian Federation in 1997

# Recommendations for young professionals

- ▶ Patriotism, morality and love of their work
- ▶ Mathematician-cryptographer should be able to build mathematical models of applied problems and show real audacity in problems formulation and solution
- ▶ A young specialist should understand that knowledge gained by him during his studies is far from sufficient to solve practical problems of information security, this is only a basis for subsequent great work in studying modern achievements in the relevant science field
- ▶ A lifetime is needed for studying new directions in science!